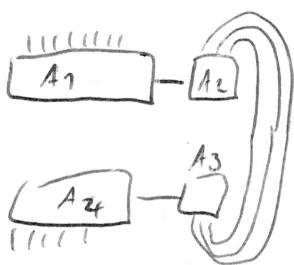


$g_A$  is a product of 4 matrices  $A_1: d^{N_A} \times D$



$$A_2: D \times d^{N_B}$$

$$A_3: d^{N_B} \times D$$

$$A_4: D \times d^{N_A}$$

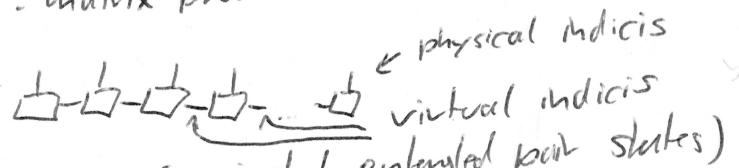
~~bond dimension~~

$$\text{rank}(g_A) = \text{rank}(A_1 A_2 A_3 A_4) \leq \min\{\text{rank } A_1, \text{rank } A_2, \text{rank } A_3, \text{rank } A_4\} \\ = D$$

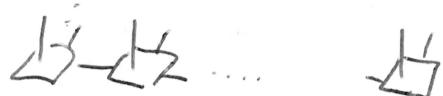
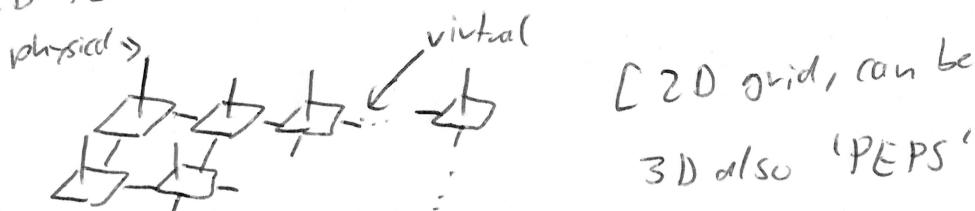
$S \leq \log \text{rank } D$ : bond dimension puts an upper bound on entanglement.  
The higher the bond dimension, the more entanglement is possible.

### Tensor network states

- 1D: matrix product states (MPS)



- 2D PEPS (projected entangled pair states)



- various other tensor network states

Remarks

$$H = \sum_{\text{local}} \epsilon_i$$

E

$$\boxed{\Delta} \approx \Delta, \Delta \text{ finite if } N \rightarrow \infty$$

(system size)

1. Ground states of local gapped Hamiltonians are well described by tensor network states.

2. PEPS (1D, 2D, 3D) fulfill an entanglement area law.

3. Many known models have exact, known tensor network ground states, for example the toric code.

4. Entanglement can be often calculated nicely with tensor network states.

→ toric code      characterized by      next lecture  
 5. Topological order is not strongly entangled, but by long-range entanglement.  
 ↓

[The reduced density matrices are not very entropic, e.g. 7% when compared to thermal states!]

6. Topological order is indicated by a topological correction to the area law:

$$S_A = | \partial A | c - \gamma + \dots$$

$\uparrow$  boundary of  $A$        $\uparrow$  terms that vanish for infinite system size  
 $\uparrow$  topological correction (less entangled)

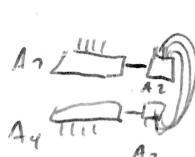


2. PEPS area law:  $|A| = N_A$ : size of  $A$ ,  $|\partial A|$ : size of boundary

$$|\psi\rangle = \boxed{\square - \square - \square - \square} \otimes \boxed{\square - \square - \square - \square} \quad \text{and } S_A = \text{tr}_B |\psi\rangle \langle \psi| ; \text{rank } S_A = ?$$



the same consideration as before?



matrices:  $A_1: d^{|\partial A|} \times D^{|\partial A|}$   
 $A_2: D^{|\partial A|} \times d^{|\partial A|}$   
 $A_3: d^{|\partial A|} \times D^{|\partial A|}$   
 $A_4: D^{|\partial A|} \times d^{|\partial A|}$

/g

$$\text{rank } S_A \leq \min(d^{|\mathcal{A}|}, d^{|\bar{\mathcal{A}}|}, D^{|\mathcal{B}A|}) = D^{|\mathcal{B}A|}$$

$$D^{|\mathcal{B}A|} < d^{|\mathcal{A}|} \quad \text{when } |\mathcal{A}| \text{ is large enough}$$

boundary  
(length)  
 $\sim L$ 
 interior  
(area)  
 $\sim L^2$

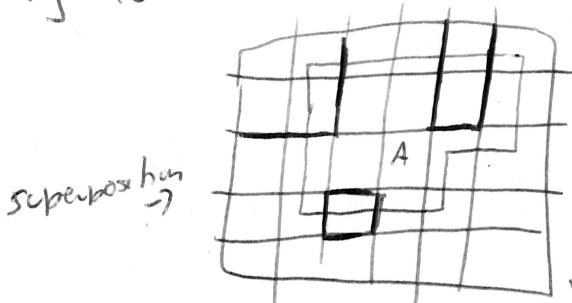
$$\Rightarrow S_0 \leq \log \text{rank}(S_A) \leq \log D^{|\mathcal{B}A|} = |\mathcal{B}A| \log D$$

area law,  $\log D$  is the constant.

Remark: generic PEPS saturate this bound. [Non-trivial statement].

Finally: How can a topological correction emerge?

- 1) Toric code - elementary (easy, idea is clear)
- 2) Toric code PEPS ('deeper', new tool: tensor network, easy to generalize)
- 3) How can  $\gamma$  be interpreted?
- 4] Toric code ground state  $|\Psi\rangle = \sum_{ij} |\Psi_i\rangle_A |\Psi_j\rangle_B = \sum_e |\Psi_e\rangle_A |\Psi_e\rangle_B$



'cut through spins', think of

loop soup

e.g. all configurations on the boundary

→ strings cross boundary an even number of times

rank  $S_A$  is limited by the number of possible boundary configurations

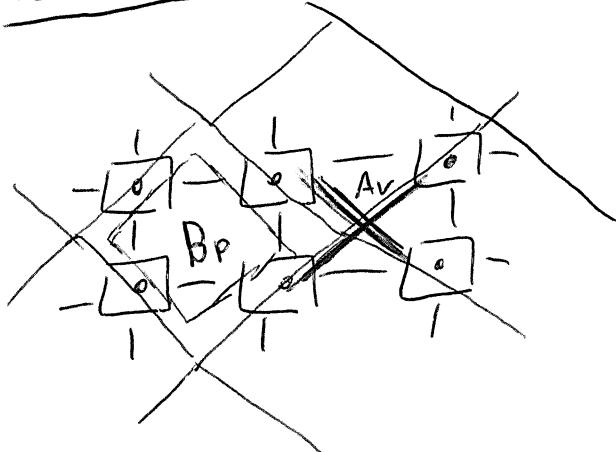
OR: 'count' states [entropic]

$$\text{How many states? } 2^{|\mathcal{B}A|} = 2^{|\mathcal{B}A|-1}$$

$$S = \underbrace{|\mathcal{B}A| \log 2}_{\text{area law}} - \underbrace{\log 2}_{\text{top. correction}}$$

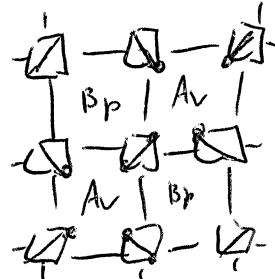
2  $\Leftarrow$  only configurations with an even number of 'strings'/spins-up-states at the boundary

## Toric Code PEPS



checkerboard  
of  $B_p$  and  $A_v$

~7



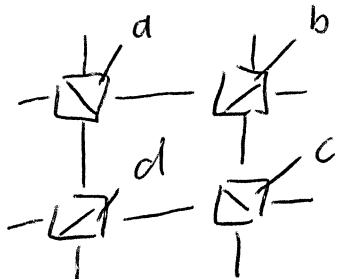
$$A_v |GS\rangle = |GS\rangle, \quad B_p |GS\rangle = |GS\rangle$$

1) parity constraint

2) superposition ↓  
of 'loop' + 'no loop'

encode 1) and 2) locally

1)



need to ensure that

$$\hat{z}\hat{z}\hat{z}\hat{z} |abcd\rangle = |abcd\rangle$$

think of parity as  $(\mathbb{Z}_2, +)$

$$\rightarrow a+b+c+d \pmod{2} = 0$$

$$d=0, 1 \quad * \quad *$$

An easy way to achieve even parity:

~~checkmark~~

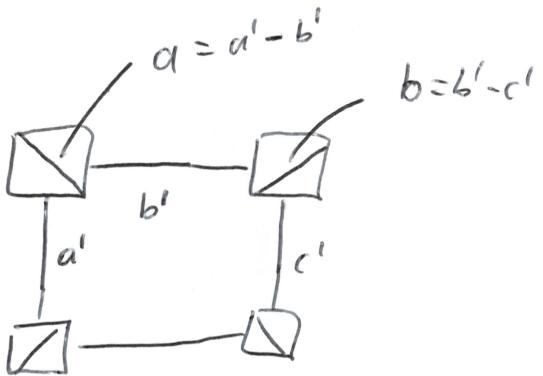
$$\begin{aligned} a &= a' - b' \\ b &= b' - c' \\ c &= c' - d' \\ d &= d' - a' \end{aligned}$$

}  $\psi_{2d}$

'define' spin states as differences  
of virtual indices

$$\Rightarrow a+b+c+d = 0$$

TC1



What about the other two indices?

$$\begin{array}{c} e' \\ \hline a' \end{array} \quad \begin{array}{c} (2) \\ a = a' - b' \end{array} \quad \leftarrow \text{to satisfy vertex constraint} \quad (1)$$

$$\begin{array}{c} f' \\ \hline b' \end{array} \quad \begin{array}{c} a = e' - f' \\ (1) \end{array} \quad " \quad (2)$$

$$\Rightarrow a' - b' = e' - f'$$

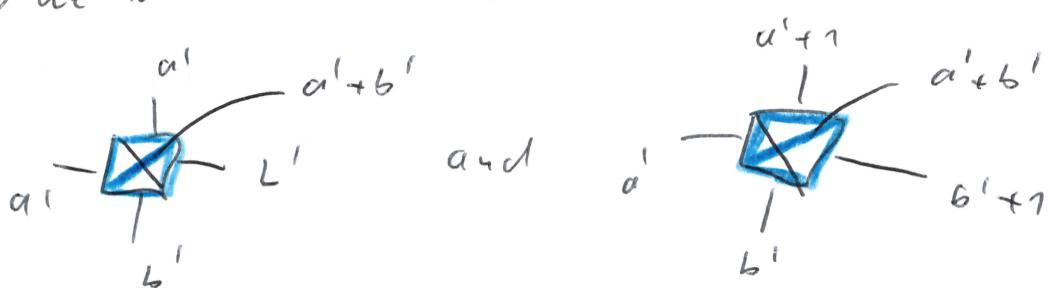
two options

$$\begin{array}{cccc|c} a' & b' & e' & f' & a' - b' = e' - f' \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 & 1 \\ \vdots & & & & \end{array}$$

either  $(a', b')$  is equal to  $(e', f')$

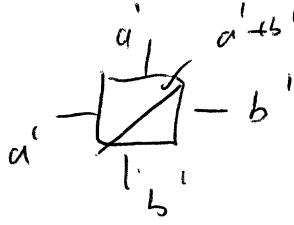
$$\text{or } e' = a' + 1 \quad f' = b' + 1$$

$\Rightarrow$  we have found the good tensors

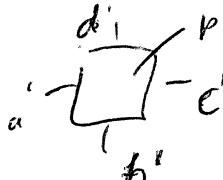


Short hand notation for  $T_{a'b'e'f'}^p = \begin{cases} 1 & \text{if } a' = e', b' = f', p = a' + b' \\ 1 & \text{if } a' = e' + 1, b' = f' + 1, p = a' + b' \\ 0 & \text{otherwise} \end{cases}$  /772

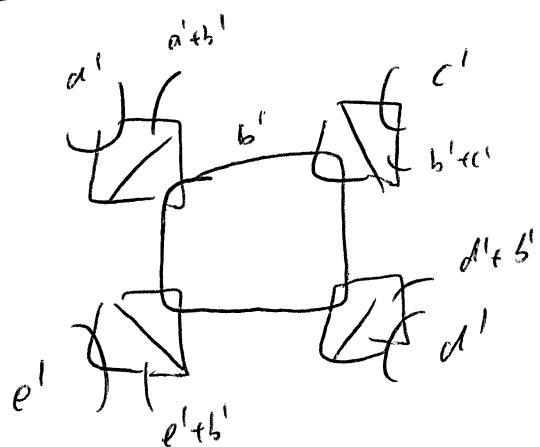
The 'first one' is sufficient.

Take 

$$T_{a'b'cd'}^p = \begin{cases} 1 & \text{if } a' = d', b' = c' \\ 0 & \text{otherwise} \end{cases}$$



show:  $B_p(4) = \frac{1}{16}$



$b'$ : closed index (are  
→ summation

Summing over  $b'$  generates  
superposition of

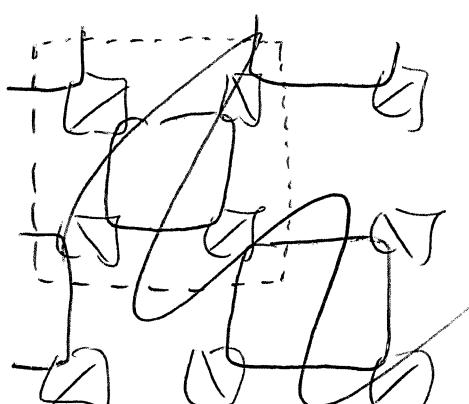
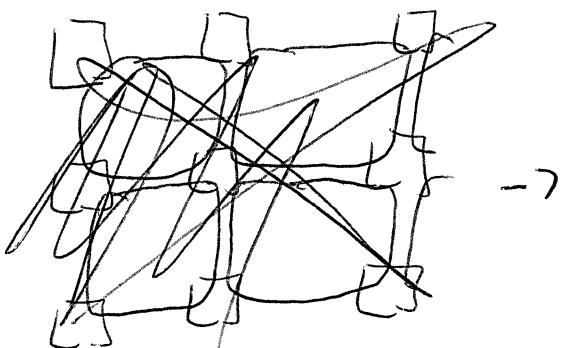
$$\frac{1+xxx}{2}$$

Here:  $b' = 0 \rightarrow |a' c' d' e'\rangle$

$b' = 1 \rightarrow |a'+1, c'+1, d'+1, e'+1\rangle$

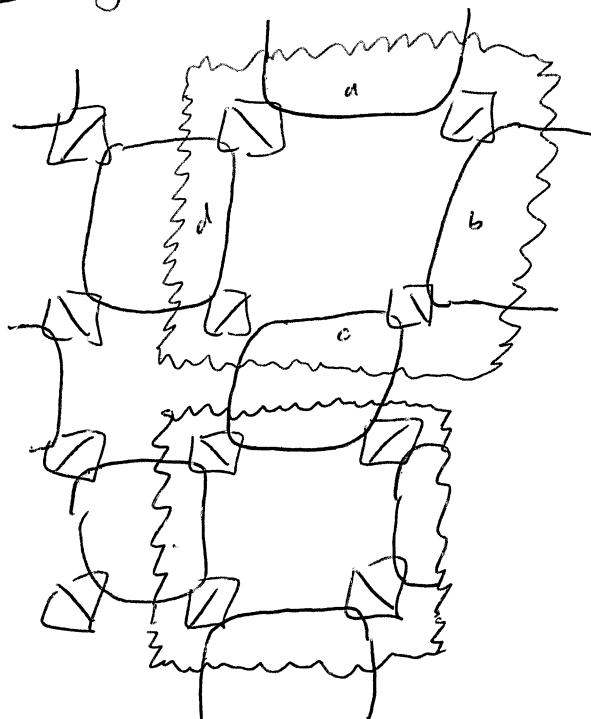
effective spin flip on all four spins

Basis

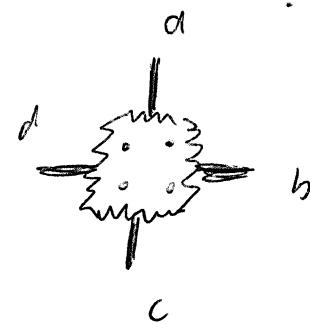


TC3

## Blocking



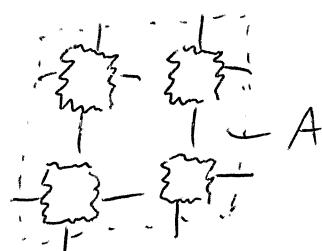
$\leadsto$



reduce redundancy and  
obtain translation invariance

## entanglement entropy

$$\text{rank } S_A \leq 2^{|A|}$$



~~but~~ actually, there is  
a symmetry in these  
tensors

$$-\boxed{x} - \underbrace{\text{---}}_{\text{Pauli-X}} - \boxed{x} = -\underbrace{\text{---}}_{\text{Pauli-X}} \left[ \begin{array}{c} \text{even} \\ \text{---} \\ \text{---} \end{array} \right] = -\text{---} - \boxed{x}$$

Every path has this symmetry!

$$x - \boxed{x} - \boxed{x} - x = -\boxed{x} - \boxed{x} - x - x = -$$

TC4

$$\text{rank } S_A = 2^{\frac{|A|}{2}} \uparrow$$

is two-to-one mapping

$$g \text{ has a flat spectrum} \rightarrow S_A = \log \text{rank } S_A$$

$$S_A = |A| \log 2 - \underbrace{\log}_f 2$$

$f$ : top correction

This is a large detour for a simple result, but the considerations generalize very straightforwardly for

$G = \mathbb{Z} \rightarrow G$ : any finite group!

$$\log 2 = \log |\mathbb{Z}_2| \rightarrow \log |G|$$

$$\text{parts} \rightarrow g_1 \cdot g_2 \cdot g_3 \cdot g_4 = e$$

$$\text{'loop sum'} \rightarrow g_1, h_1, i_1, k_1 \xrightarrow[X]{} g_{ox}, h_{ox}, i_{ox}, k_{ox}$$

Kitaev group quantum double models

$\rightarrow$  include non-Abelian anyon models, e.g.

$$D(S_3)$$